

Belief propagation decoding of quantum channels by passing quantum messages

arXiv:1607.04833

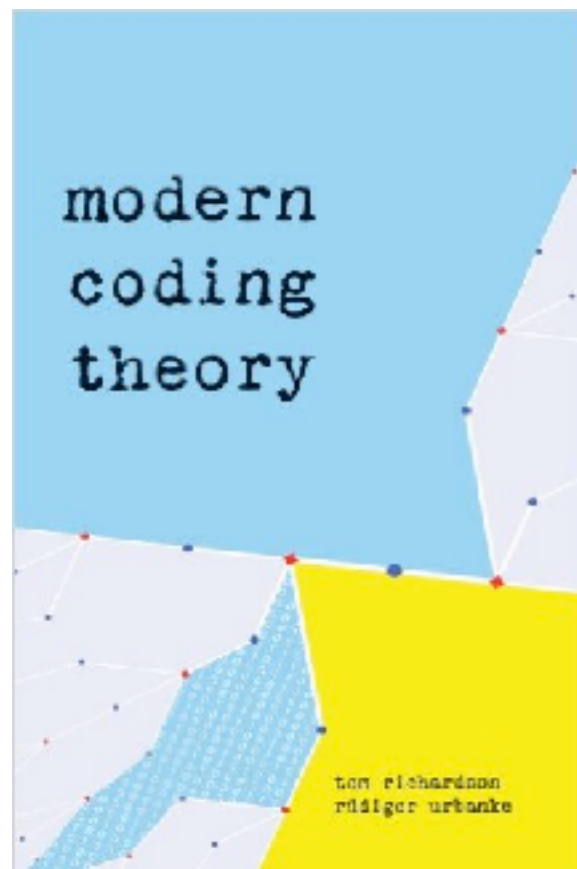
QIP 2017

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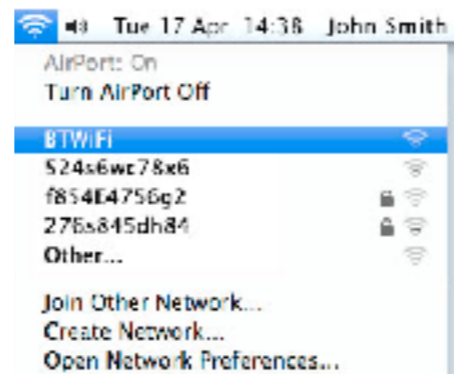
Joseph M. Renes
ETH zürich



To do research in quantum information theory, pick a favorite text on classical information theory, open to a chapter, and translate the contents into quantum-mechanical language.



Based on belief propagation decoding



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Spatially Coupled Ensembles Universally Achieve Capacity Under Belief Propagation

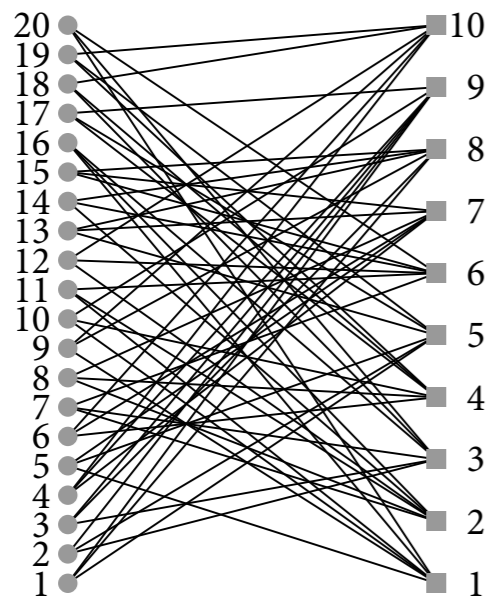
Shrinivas Kudekar, Tom Richardson, *Fellow, IEEE*, and Rüdiger L. Urbanke

Abstract—We investigate spatially coupled code ensembles. For transmission over the binary erasure channel, it was recently shown that spatial coupling increases the *belief propagation* threshold of the ensemble to essentially the *maximum a priori* threshold of the underlying component ensemble. This explains why convolutional LDPC ensembles, originally introduced by Felström and Zigangirov, perform so well over this channel. We show that the equivalent result holds true for transmission over general binary-input memoryless output-symmetric channels. More precisely, given a desired error probability and a gap to capacity, we can construct a spatially coupled ensemble that fulfills these constraints *universally* on this class of channels under belief propagation decoding. In fact, most codes in this ensemble have this property. The quantifier *universal* refers to the *single* ensemble/code that is good for *all* channels but we assume that the channel is known at the receiver. The key technical result is a proof that, under belief-propagation decoding, spatially coupled

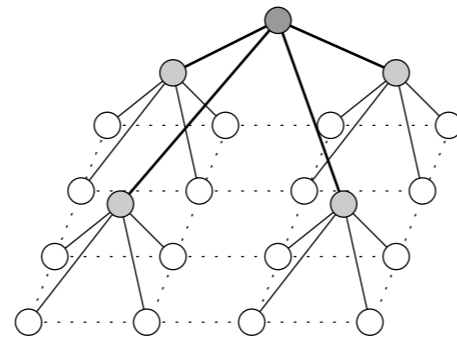
In the first 50 years, coding theory focused on the construction of *algebraic* coding schemes and algorithms that were capable of exploiting the algebraic structure. Two early highlights of this line of research were the introduction of the Bose–Chaudhuri–Hocquenghem (BCH) codes [5], [6] as well as the Reed–Solomon (RS) codes [7]. Berlekamp devised an efficient decoding algorithm [8], and this algorithm was then interpreted by Massey as an algorithm for finding the shortest feedback-shift register that generates a given sequence [9]. More recently, Sudan introduced a list decoding algorithm for RS codes that decodes beyond the guaranteed error-correcting radius [10]. Guruswami and Sudan improved upon this algorithm [11] and Koetter and Vardy showed how to handle soft information [12]. Another important branch started with the introduction of convolutional codes [13] by Elias and the introduction of the

Quantum belief propagation decoding?

Belief propagation: **message passing** algorithm for performing **inference** in a **graphical model**



decoder infers channel input from output;
code can be described by a graphical model



early precursor: Bethe-Peierls
approximation in statistical physics

many applications in statistics and machine learning besides coding:
inference, optimization, constraint satisfaction

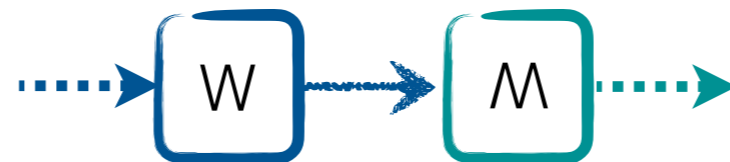
Quantum belief propagation decoding?

- Can use classical algorithm for usual stabilizer decoding

- Optimal for: 

- But not, e.g. amplitude damping 

Let's consider CQ channels for simplicity



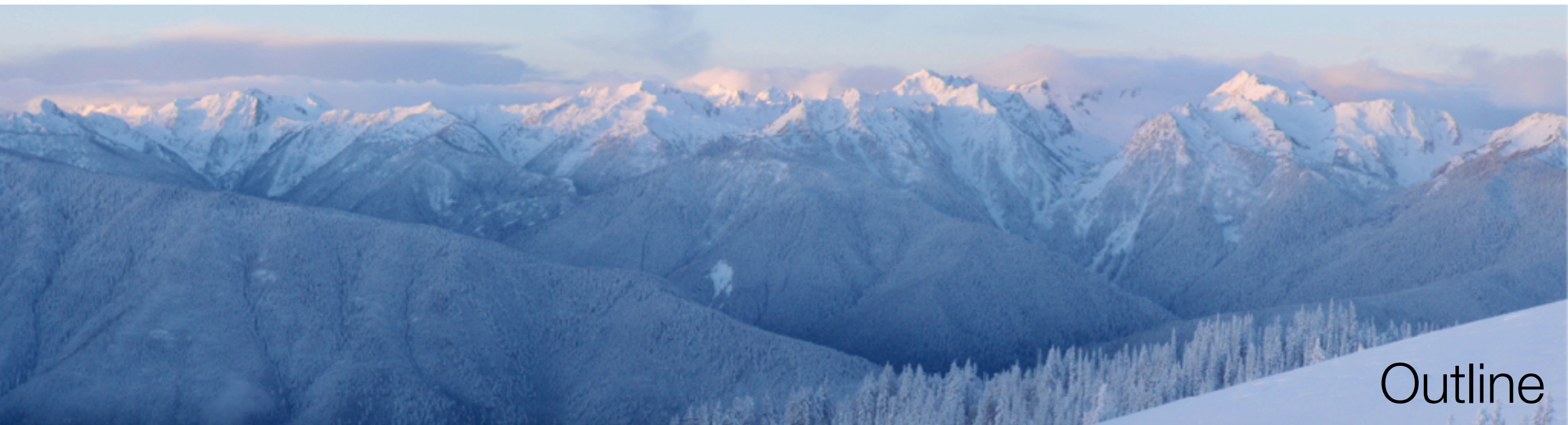
Need to infer channel input from quantum output
(not trying to compute marginals of quantum states)

Results

arXiv:1607.04833

- BP decoder for pure state channel, tree codes
- Also works for polar codes:
Efficient, capacity-achieving decoder for BPSK over lossy Bosonic channel
- And for quantum communication, part of conjugate basis decoder:
Efficient, capacity-achieving decoder for amplitude damping

1. Coding setup
2. Factor graphs
3. Classical BP

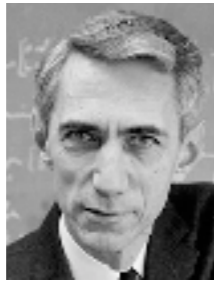


Outline

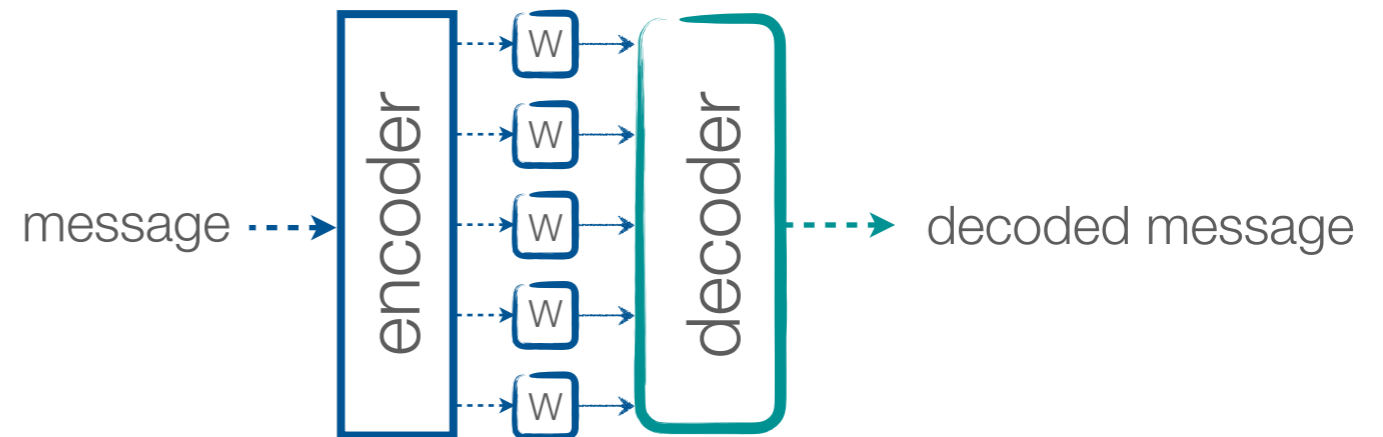
National Park Service

4. Quantum BP
5. Many open questions

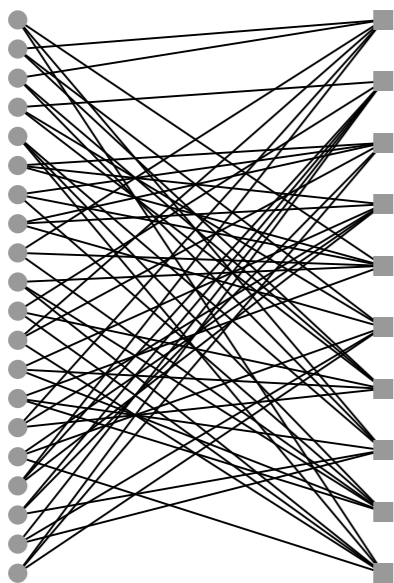
Coding setup



Shannon scenario: stochastic iid noise.
not adversarial; fault-free decoding



Linear code



bitwise decoding of random codeword

classically, marginalize joint distribution

$$P_{X_1^n Y_1^n = y_1^n} \rightarrow P_{X_i Y_1^n = y_1^n}$$

quantumly, perform Helstrom
measurement for each bit

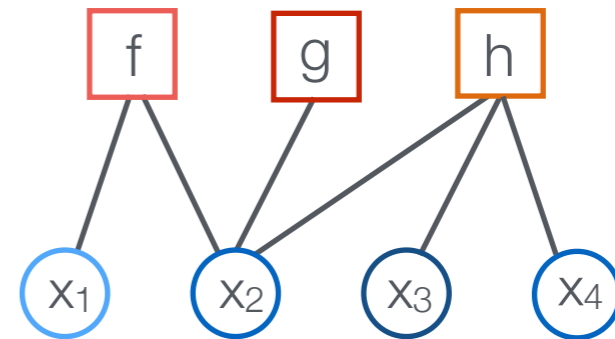


$$\text{not } \rho_{X_1^n B_1^n} \rightarrow \rho_{X_i B_1^n}$$

Factor graphs

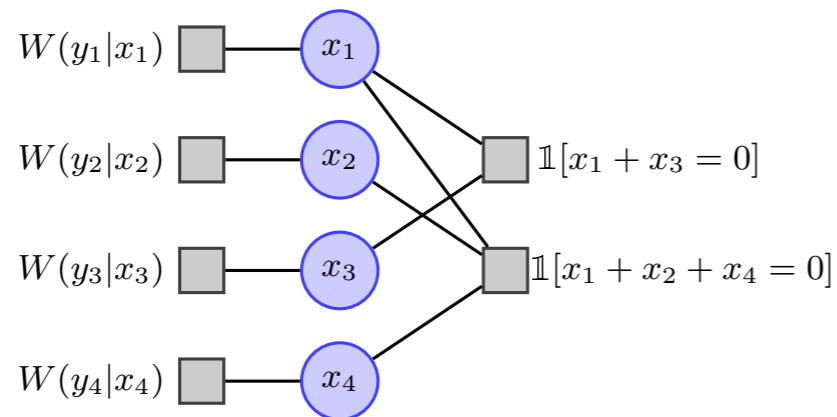
factorizable joint probability

$$P(x_1, x_2, x_3, x_4) = \frac{1}{Z} f(x_1, x_2) g(x_2) h(x_2, x_3, x_4)$$



easier to compute marginals:

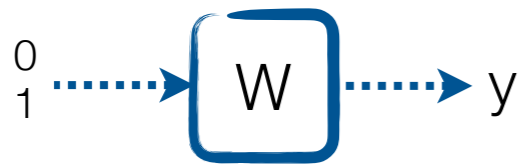
$$P(x_1) = \sum_{x_2, x_3, x_4} P(x_1, x_2, x_3, x_4) = \frac{1}{Z} \sum_{x_2} f(x_1, x_2) g(x_2) \left(\sum_{x_3, x_4} h(x_2, x_3, x_4) \right)$$



channel input & output distribution:

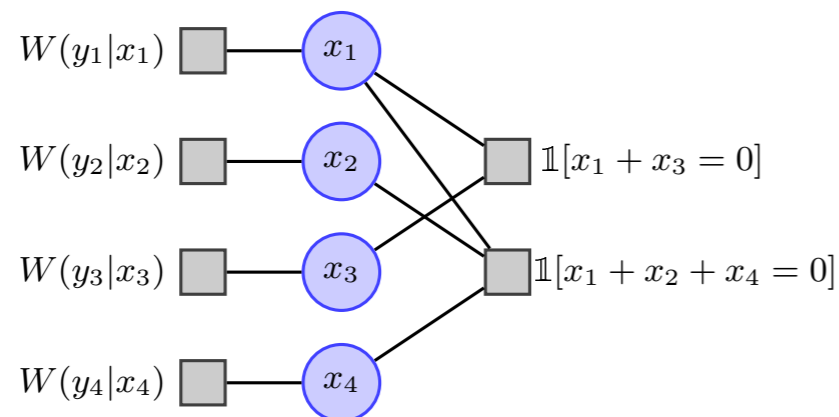
$$P_{X^n Y^n}(x_1^n, y_1^n) = \frac{1}{|C|} \mathbb{1}[x_1^n \in C] \prod_{j=1}^n W(y_j | x_j).$$

Classical BP decoding



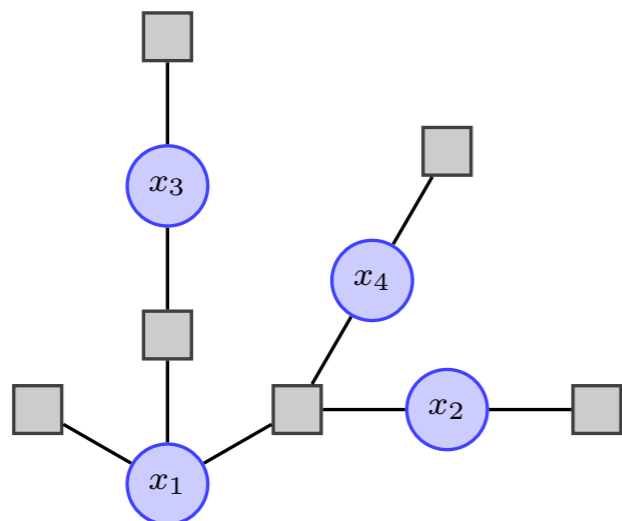
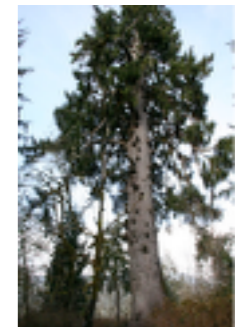
symmetric binary input channel:
only care about likelihood ratio

$$\ell(y) = \frac{W(y|1)}{W(y|0)}$$



$x_i \rightarrow y_1^n$ is a channel;
want $\ell(y_1^n)$

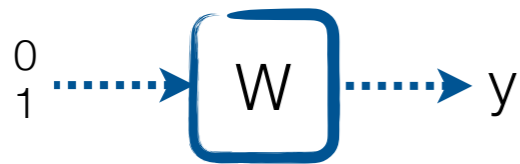
assume tree factor graph



for each node,
associate a channel to all its leaves

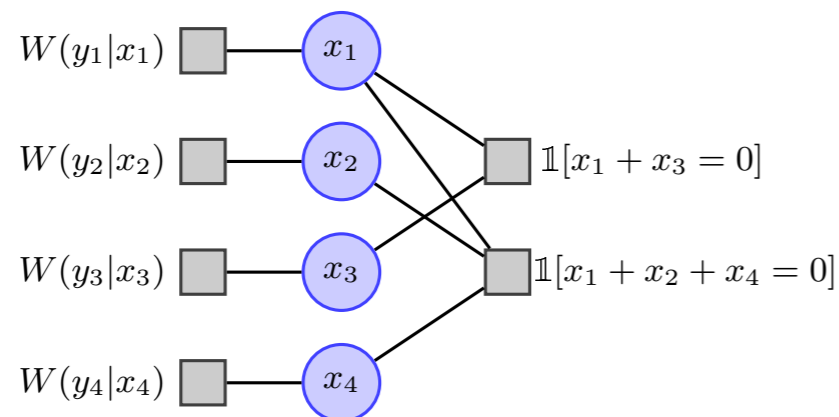
BP recursively computes $\ell(y_1^n)$
starting from the leaves

Classical BP decoding



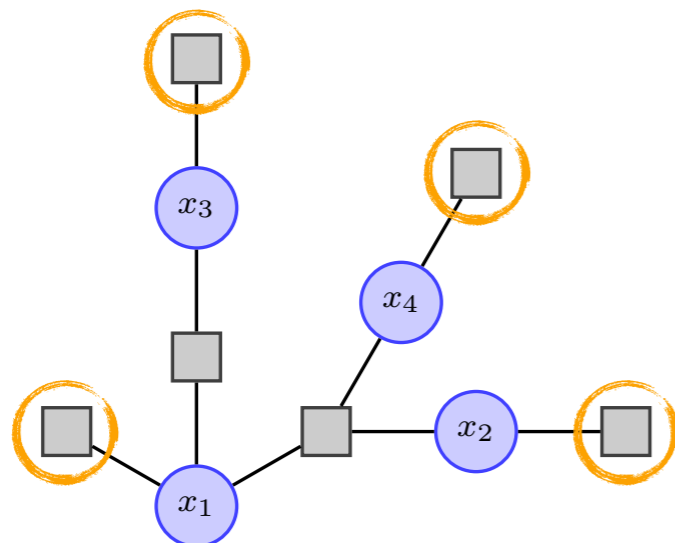
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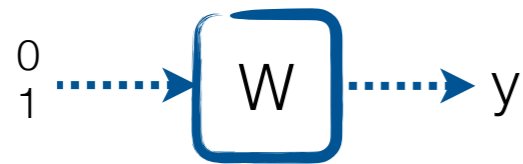
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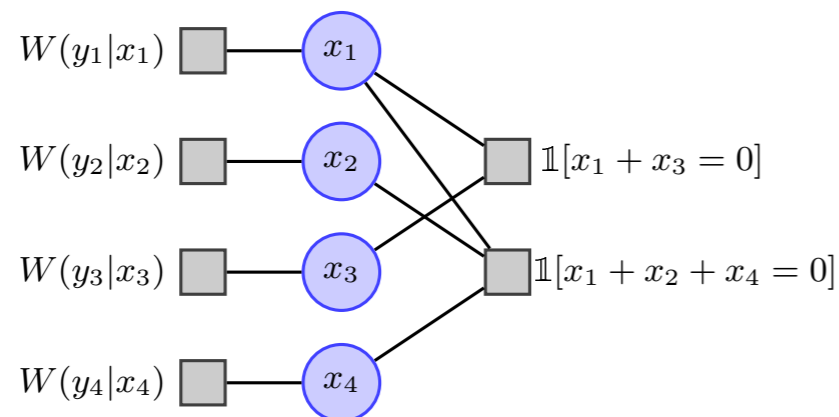
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Classical BP decoding



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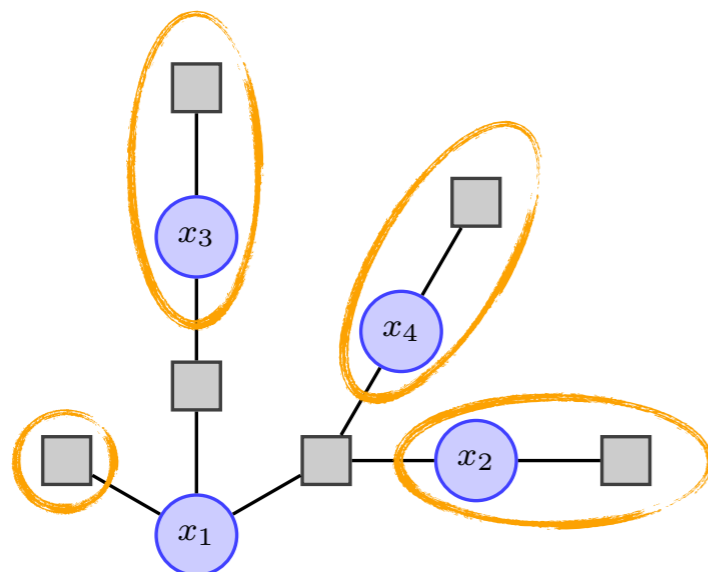
$$\ell(y) = \frac{W(y|1)}{W(y|0)}$$



$x_i \rightarrow y_1^n$ is a channel;
want $\ell(y_1^n)$



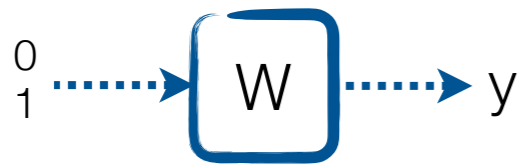
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for each node,
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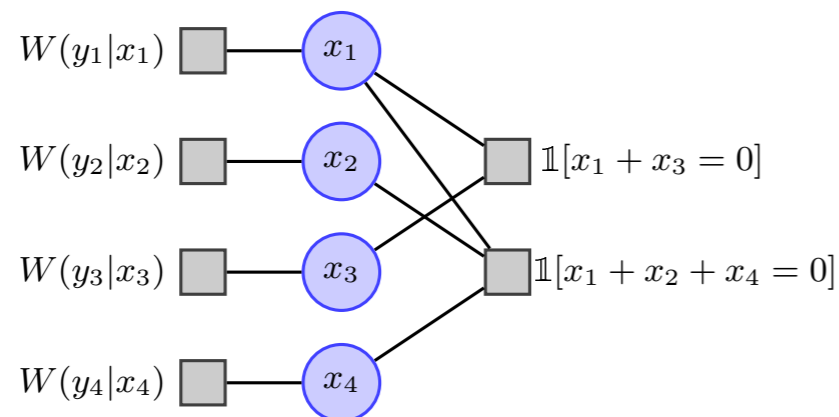
BP recursively computes $\ell(y_1^n)$
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Classical BP decoding



symmetric binary input channel:
only care about likelihood ratio

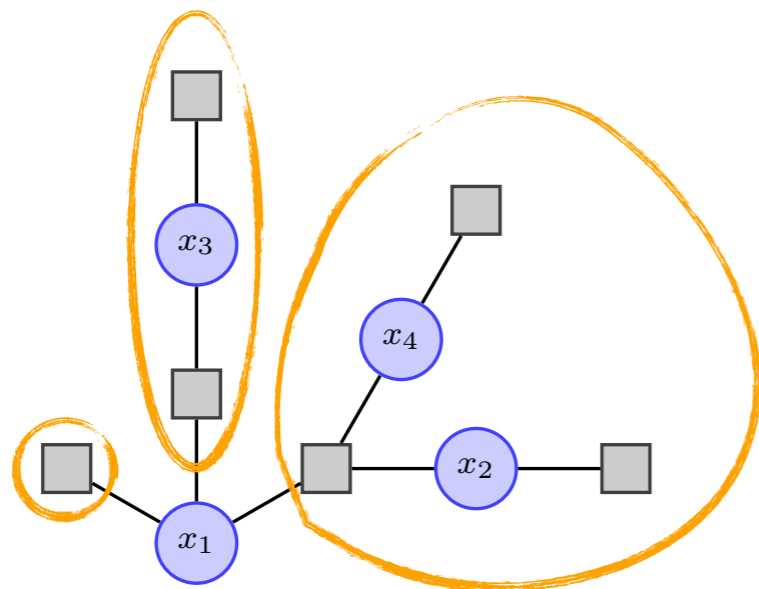
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$x_i \rightarrow y_1^n$ is a channel;
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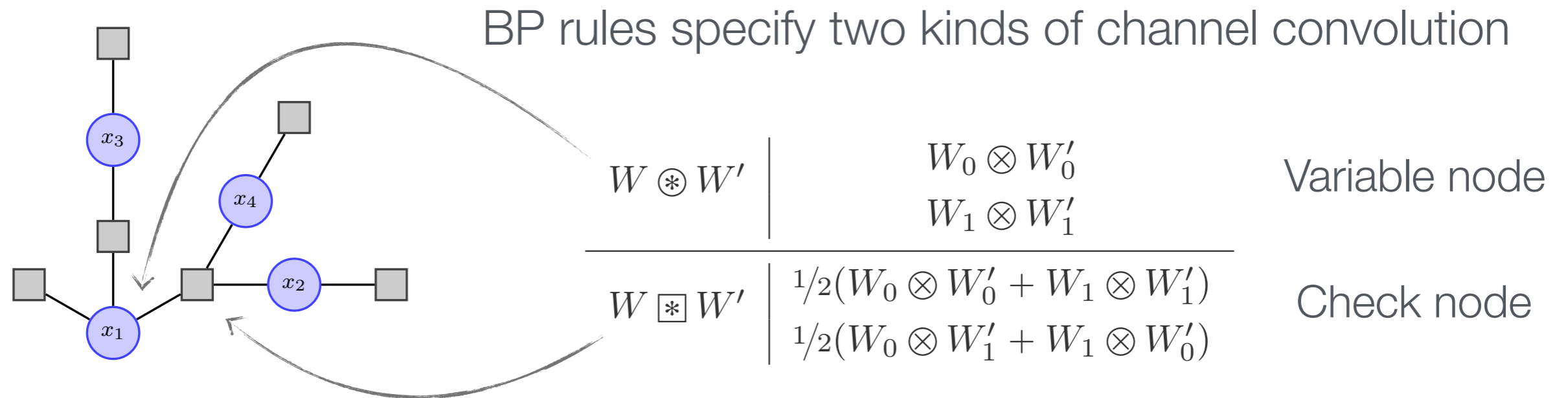
assume tree factor graph



for each node,
associate a channel to all its leaves

BP recursively computes $\ell(y_1^n)$
starting from the leaves

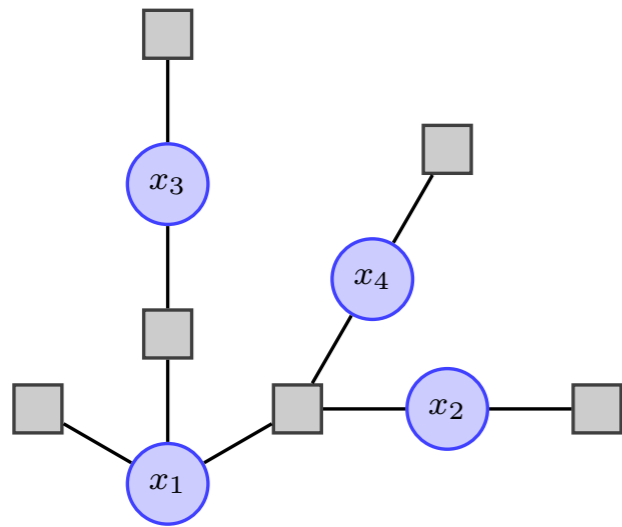
Classical BP decoding



- BP finds exact marginals on trees
- Can run algorithm for all codeword bits concurrently
- Also works on loopy LDPC factor graphs



Quantum BP decoding



Now the outputs are quantum, no likelihood function

Want to perform Helstrom measurement

Can we combine Helstrom measurements along the branches?

Can we “repackage” the outputs as with likelihood?

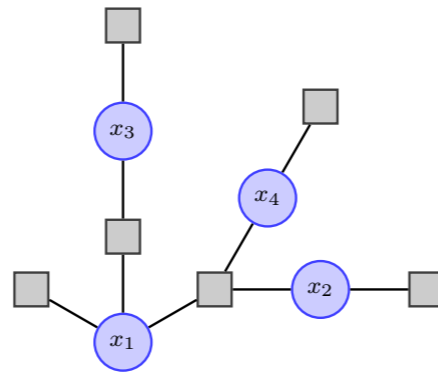
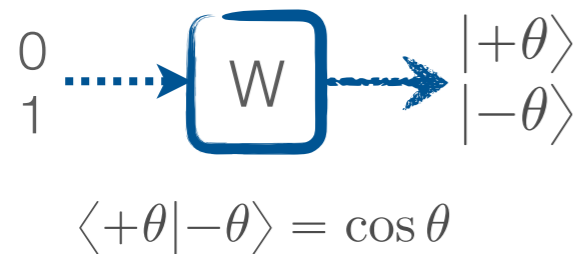
Try the simplest case:
pure state outputs



$$\langle +\theta | -\theta \rangle = \cos \theta$$

optimal measurement is σ_x

Quantum BP decoding



$W \circledast W'$	$W_0 \otimes W'_0$ $W_1 \otimes W'_1$
$W \boxtimes W'$	$\frac{1}{2}(W_0 \otimes W'_0 + W_1 \otimes W'_1)$ $\frac{1}{2}(W_0 \otimes W'_1 + W_1 \otimes W'_0)$

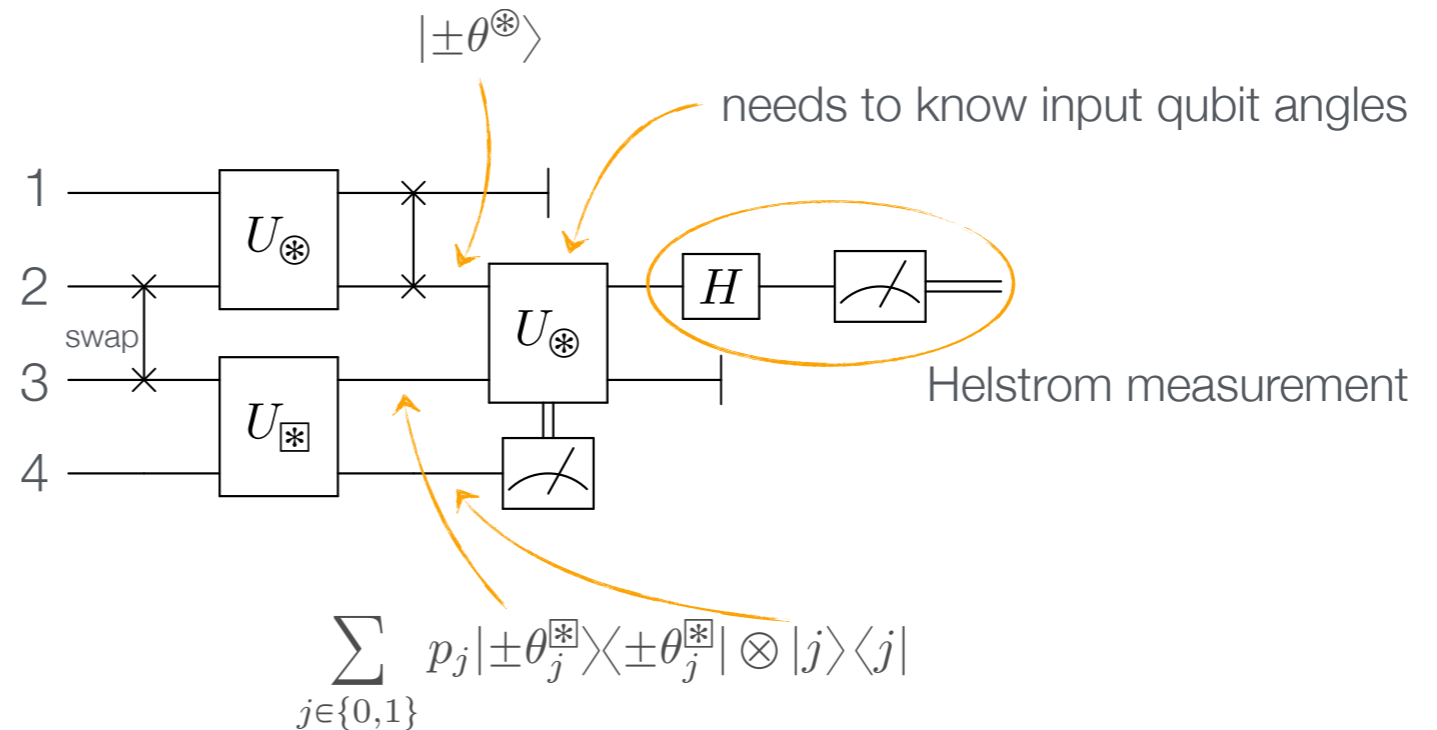
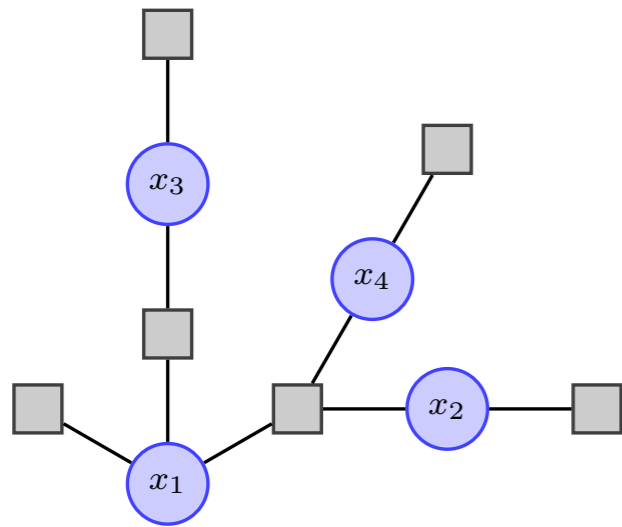
\circledast convolution yields pure states: $|+\theta\rangle \otimes |+\theta'\rangle$ $|-\theta\rangle \otimes |-\theta'\rangle$ \longrightarrow $|\pm\theta^{\circledast}\rangle$ $\cos \theta^{\circledast} = \cos \theta \cos \theta'$ repackage in a single qubit with $U_{\circledast}(\theta, \theta')$

\boxtimes convolution gives a *heralded* mixture of pure states!

unitary U_{\boxtimes} gives $\sum_{j \in \{0,1\}} p_j |j\rangle \langle j| \otimes |\pm\theta_j^{\boxtimes}\rangle \langle \pm\theta_j^{\boxtimes}|$

$$p_0 = \frac{1}{2}(1 + \cos \theta \cos \theta') \quad \cos \theta_0^{\boxtimes} = \frac{\cos \theta + \cos \theta'}{1 + \cos \theta \cos \theta'} \quad \cos \theta_1^{\boxtimes} = \frac{\cos \theta - \cos \theta'}{1 - \cos \theta \cos \theta'}$$

Quantum BP decoding



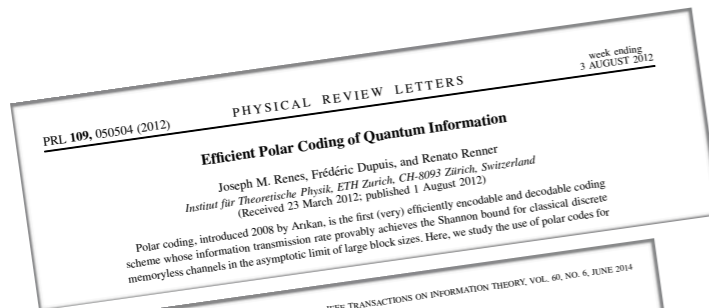
- pass qubits and some classical bits: “sufficient statistic”
- decode all codeword bits sequentially, unwinding each time
- $O(n^2)$ implementation of all Helstrom measurements

Quantum BP & polar codes



Variable and check convolutions = “better” and “worse” synth channels

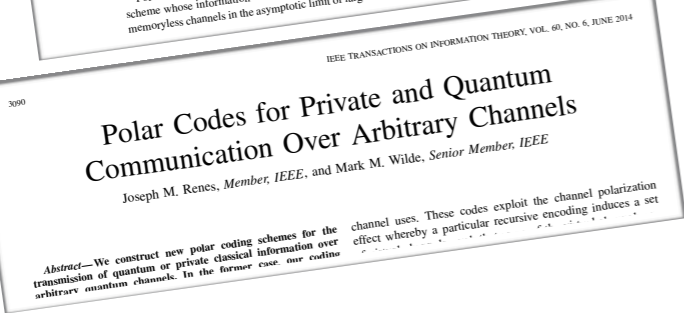
Polar decoder has tree structure (for message not codeword bits, tho)



Quantum polar decoder uses polar decoder for classical amplitude and phase info

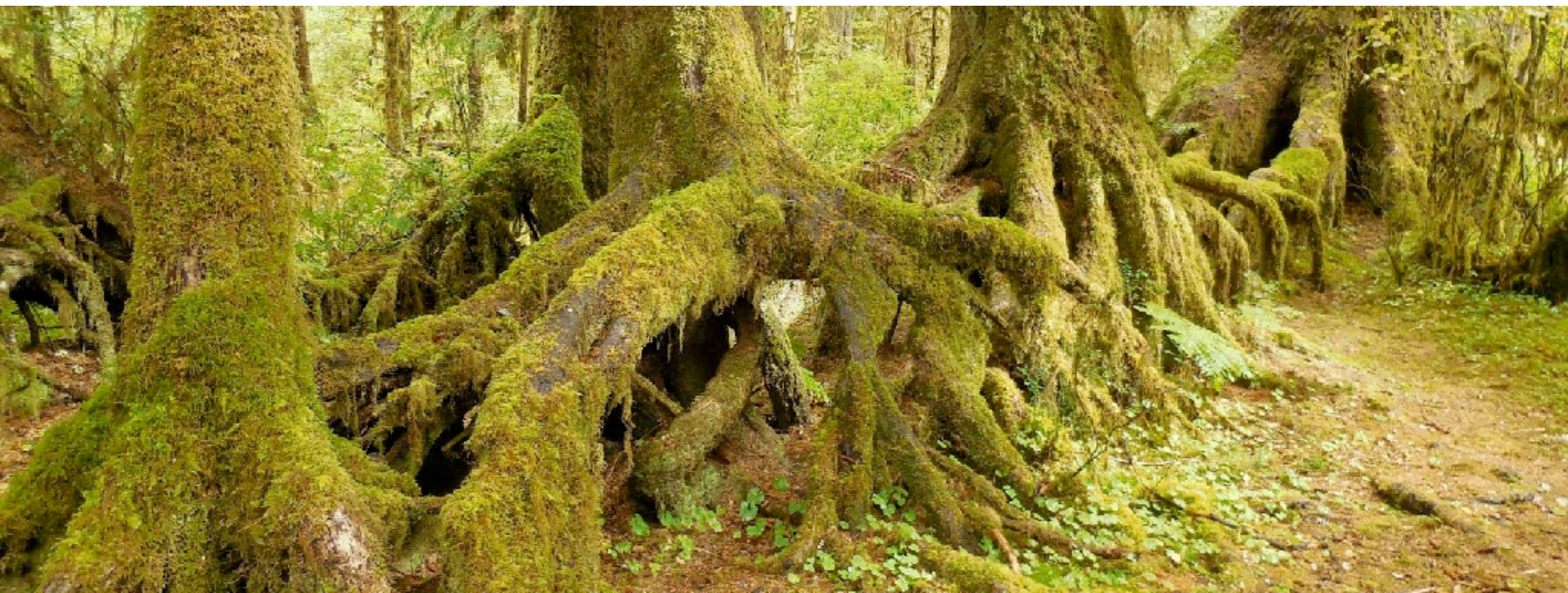


: amplitude = classical Z channel,
phase = pure state channel



$O(n^2)$ decoder for capacity-achieving quantum polar code for amplitude damping

Open



oldsloat.blogspot

Questions

Open

other channels? e.g. classical coding for amplitude damping?

loops? spatially-coupled LDPC codes?

use density matrix BP? need sufficient statistics, local tree to global

relation to tensor networks?

junction-tree algorithm?

Viterbi decoder (blockwise decoder)?

fully quantum version?

other tasks besides decoding?

Questions